

## Unit : 2 Regular Languages

### \* Regular Languages :

Regular Languages is a language that can be expressed with a Regular expression.

A Regular Languages is set of string which made up of characters.

Regular Languages is the most restricted types of Languages.

Regular Languages are expected by Finite Automata.

→ Operations of Regular Languages.

(i) Union : IF A and B are two regular languages then Union of  $A \cup B$  is also union.

(ii) Intersection : IF A and B are two regular languages then intersection of  $A \cap B$  is also follow intersection operation.



ciii) Kleen Closure: IF A is a regular language then its kleen closure  $A^*$  is also be a regular language.

\* Regular Expression:

Regular Language accepted by Finite automata can used to described by Regular Expression.

Regular Expression is used to define sequence of pattern of string.

Ex.  $(a + b)^*$

Regular Expression =  $\{ \Lambda, a, b, ab, aa, bb, \dots \}$

=> Note:

\* -> Clean closure which contain null string which is denoted by  $\Lambda$  or  $\epsilon$



+  $\rightarrow$  Positive closure which does not contain null string

$$a^* = \{\epsilon, a, aa, \dots\}$$

$$a^+ = \{a, aa, aaa, \dots\}$$

Ex. Example of Regular Expression.

a String having zero or more a

$\rightarrow$  Regular Lan. =  $\{\epsilon, a, aa, \dots\}$

Regular Exp. =  $a^*$

b String end with a

$\rightarrow$  Regular Lan. =  $\{a, ba, aa, aaa, \dots\}$

Regular Exp. =  $(a + b)^* a$

c String start with a and End with b

$\rightarrow$  R.L =  $\{ab, aab, abb, aabb, \dots\}$

R.E =  $a(a + b)^* b$

d String having one or more a

$\rightarrow$  R.L =  $\{a, aa, aaa, \dots\}$

R.E =  $a^+$



e Binary String start with 0 and end with 1

→ R.L. = {01, 001, 011, 0011, ... }  
R.E =  $0(0+1)^*1$

f String should contain at least 3 One

→ R.L. = {111, 1011, 0111, 1101, ... }  
R.E =  $(0+1)^*1(0+1)^*1(0+1)^*1(0+1)^*$

g String should have odd length

→ R.L. = {0, 1, 010, 011, 01101, ... }  
R.E =  $(0+1)((0+1)(0+1))^*$

h String end with 1 and not contain 00

→ R.L. = {1, 01, 011, 1010, ... }  
R.E =  $(1+01)^*$

i String of length 6 or less

→ R.E =  $(0+1+\wedge)^6$



j String Begins or ends with 00 or 11

→ R.L. = {0011, 00011, 00111, ...}  
R.E. =  $(00+11)(0+1)^* + (0+1)^*(00+11)$

\* Indistinguishable String:

Suppose, L is a Regular Language and X and Y is a string which present in L.

Suppose Z is one string.

If combination of XZ and YZ both are present in L.

So, this string is called Indistinguishable string.

\* Distinguishable String:

Suppose, L is a Regular Language and X and Y is a string which present in L.

Suppose Z is one string,



If combination of  $XZ$  is present in  $L$  and  $YZ$  is not present in  $L$ .

So, this string is called Distinguishable string.

Ex. Regular Language  $L = \{0, 1, 00, 11, 01, 10, \dots\}$

$\rightarrow$  Take  $X = 0$  and  $Y = 1$   
Suppose  $Z = 1$ .

For  $L$ , Distinguishable string,

$$XZ = 01 \in L$$

$$YZ = 11 \in L$$

So, this string is Indistinguishable string.

So, this string is not distinguishable string. because both string are present in  $L$ .